

# A C++/CUDA DSL for Object-oriented Programming with Structure-of-Arrays Layout



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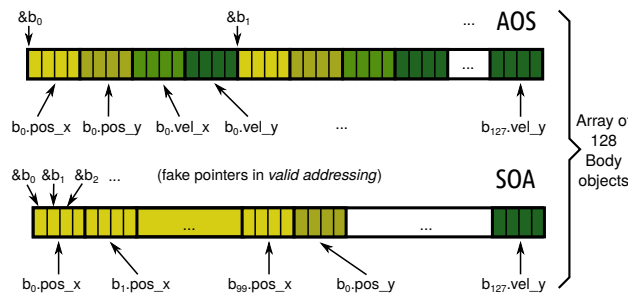
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<https://github.com/prg-titech/ikra-cpp>

**Context:** HPC with uniformly structured data (e.g., n-body simulation, traffic flow simulation)

**Goal:** SOA memory layout (good for caching, vectorization, parallelization) with C++ Notation

Pointer instead of IDs, Method Calls, **new** Keyword, Templates, Member of Object/Pointer Operator, Future work: virtual functions



**Object Creation:** `Body *p = new Body(1.0, 2.0);`  
`Body *q = Body::make(10, 1.0, 2.0);`

**Field Access:** `p->vel_x = p->vel_y = 1.5;`

**Member Functions:** `p->move(0.5);`  
`forall(&Body::move, q, 10, 0.5);`

## Related Work

Robert Strzodka. *Abstraction for AoS and SoA Layout in C++*. GPU Computing Gems Jade Edition, pp. 429-441, 2012. [ First DSL approach in C++. Supports easy change between AOS and SOA layout. Complicated notation. Potentially large mem. footprint if fields have different size. ]

Holger Homann, Francois Laenen. *SoAx: A generic C++ Structure of Arrays for handling particles in HPC code*. Comp. Phys. Comm., Vol. 224, pp. 325-332, 2018. [ Simpler notation than [Strzodka12] for single struct. Still not like standard C++. Expression Templates to avoid memory allocation for temporary results in large arith. expressions. ]

```
class Body : public SOA<Body> {
public: INITIALIZE_CLASS
float_pos_x; float_pos_y;
float_vel_x; float_vel_y;
```

```
Body(float x, float y) : pos_x(x), pos_y(y) {}
```

```
void move(float dt) {
pos_x = pos_x + vel_x * dt;
pos_y = pos_y + vel_y * dt;
}
```

```
}; GPU mode: Use DEVICE_STORAGE. Max. #objects
```

```
HOST_STORAGE(Body, 128);
```

`char buffer[128 * 16];`  
Large enough to store 128 objects (four float[128] arrays)

## "Fake" Pointers encode Object IDs

There are various encoding techniques. Need to specify both an encoder (object construction) and a decoder (field access).

a) Zero Addressing: `&obj_id = id`

```
void* Body::operator new() {
return (void*) size++;
} // encoder
```

```
int field<T, idx, offset>::id() {
Body* ptr = ((char*) this)
- idx*sizeof(field<...>);
return (int) ptr;
} // decoder
new keyword virtual functions
```

b) Valid Addressing: `&obj_id = buffer + id`

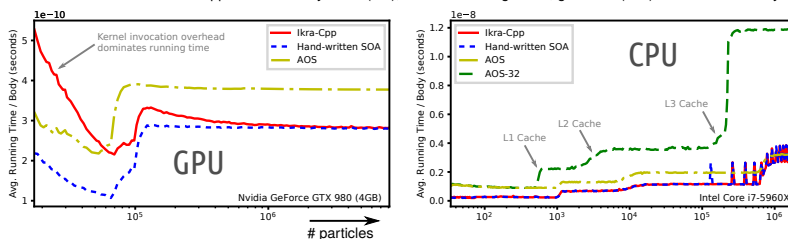
c) First Field Addressing: `&obj_id = buffer + sizeof(T) * id`

## Results & Main Insights

- ★ Field access (decoding object IDs + calculating memory addresses) is as efficient as array access in hand-written SOA code (*strided mem. access*).  
Zero Addressing: `data_ptr_vel_x(ptr) = 0x600400 + 4 * ptr`  
Valid Addressing: `data_ptr_vel_x(ptr) = -0x11FFC00 + 4 * ptr`
- ★ Main Limitation: How well can the compiler optimize this code? Experiments performed with Zero Addressing.  
gcc 5.4.0: compiler hints necessary (constexpr)  
clang 3.8, 5.0: works for simple examples, loop vectorization fails (because of single buffer array)  
Future work: Reimplement with ROSE Compiler (a powerful C++ preprocessor)

```
float test(Body* ptr) {
return ptr->vel_x;
}
```

Microbenchmark: Iterative Application of `Body::move(0.5)`, Zero Addressing Mode, gcc 5.4.0 (-O3). Identical assembly.



SO, how does it work?

SOA field types index offset

```
float_vel_x;
field<float, 2, 8> vel_x;
```

Make `field<float>` behave like a float

Implement implicit conversion and assignment operator.

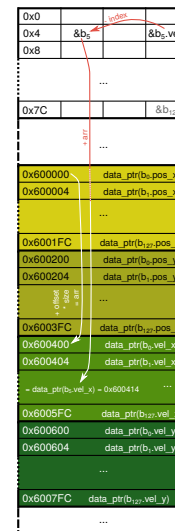
```
field<T, idx, offset>::operator T&() {
return *data_ptr();
}
```

Calculate memory location of field value

```
T* field<T, idx, offset>::data_ptr() {
T* arr = (T*) (buffer + 128*offset);
return arr + id();
}
(padding area needed for valid + first field addressing)
```

Mem. Layout Example

(Zero Addr., (Float) b<sub>i</sub>.vel\_x)



More OOP features, but harder to optimize