

A Layer-based Approach to Hierarchical Dynamically-scoped Open Classes

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Introduction



- "Open Classes" in Ruby
 - Modify existing classes or modules
 - Add or overwrite methods
- Why Open Classes?
 - Object-oriented auxiliary methods
 e.g.: 5.minutes + 9.hours
 - Multi-dimensional separation of concerns [Tarr99]
 - Bug fixing (*monkey patching*)
- Support in programming languages
 - Ruby: open classes
 - Smalltalk: extension methods
 - Python: modifiable method dictionary

Introduction



Open Classes in Ruby by example

in standard library

class Fixnum

end.

in a different component

class	Fixnum
	def minutes
	return self * 60
	end
	def hours
	return self * 3600
	end
end	



Example 1: WebPage Library [Takeshita13]



- A library: WebPage renders HTML and might show popups
- Two applications: using WebPage
 - Browser: should *not* show popups
 - Viewer: should show popup



The Problem: Global Visibility





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The Problem: Global Visibility





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Example 1: With Open Classes





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The Problem with Open Classes



- Global Visibility
 - Modifications are visible everywhere
 - Other components (e.g., Viewer) can break
 → "Destructive Modifications"
- Solution: Locality of Changes



Idea: scope control of modifications

- Using only classes (vs. classboxes, method shells etc.)
- Reusability through Ruby *modules* (or mixins) (vs. new syntax for refinements)
- Consistent with Ruby's language features: take into account class nesting hierarchy
- Amenable to other programming languages with
 - object-based, class-based
 - unit of reuse (e.g., mixins/modules, traits, ...)
 - (class nesting hierarchy)

Extension Classes



- Modifications are defined as "inner classes"
- only visible from "enclosing classes" (details follow)



• directly?

- from a different class?
- via a different class?
- via another method in the modified class?
- via a sibling inner class?
- via a superclass?
- via a subclass?

Subtleties of visibility

should modification visible when it is called:





in the context

Modifications are visible

Our principle of visibility

- of an enclosing class, and
- as long as the context remains within enclosing/ sibling classes (cf. COP)





Activation Rule



- Set of active classes S = { }
- When calling a method **C.foo**: Add C to S

browser.open(...) # S += Browser

```
viewer.check(...) # S += Viewer
```

Deactivation Rule



- Restore original S when returning from a method call
- When calling C.foo, deactivate all classes a ∈ S, where C ∉ scope(a)
 - Intuitively: scope(C) is a set of classes that are compatible with the modifications defined by C
 - Mathematically: scope(C) = { C } U all target classes
 - Definition will be extended later

class Browser
 partial
 class ::WebPage; end

end

scope(Browser) =
{ Browser, WebPage }

- Browser is compatible with Browser's modifications
- WebPage is compatible with Browser's modifications

Example 1: Overview





class Application
 def main
 Browser.new.open("http://www.titech.ac.jp")
 Viewer.new.check("secret.html")
 end

end



Object {Object (, WebPage, Browser, Viewer, Application)}





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Object {Object (, WebPage, Browser, Viewer, Application)}







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Object {Object (, WebPage, Browser, Viewer, Application)}



S = { Browser }











Object (Object (, WebPage, Browser, Viewer, Application)) S = { Browser, WebPage } Application WebPage Browser L | {Browser, WebPage} Viewer} WebPage.popup WebPage main popup check open popup do nothing open WebPage.open Browser.open Application.run ►



S = { Browser }





S = { Application }





 $S = \{ Viewer \}$











Example 1: Variations





Reusability with Modules



- Classes and modules can define partial classes
- Modifications are active in including classes (as if they were defined there directly)

module NoPopup	class Browser	
partial	include NoPopup	
	#	
class ::WebPage	end	
def popup		
#	class Viewer	
end	include NoPopup	
end	#	
end	end	

Class Activation Schemes



 How can we ensure that a class M is active when running code from class C?

Class-based Activation

- Control flow passes through class M
- For every class c that is visited on the way to C: $c \in scope(M)$
- M pushes modifications to C (cf. local rebinding/ dynamic scoping)

M is a module/mixin C includes M

C requests modifications from M

Mixin-based Activation

include NoPopup

end TiTech / HPI

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Hierarchical Scoping



- How do we share modifications among an entire class nesting hierarchy?
- Modifications of class C should affect all classes that are nested inside C

Example 2: Overview



Object {Object, AddressBook, AddressBook.Address, Networking, Networking.Address, String}

AddressBook, Address, String}	Networking {Networking, Address, String}	String		
- Address ^{[{Address}} to_address	- Address			
L String	– String			
└ to_address	└ to_address			
– Pinging				
	└ ping			

Activation / Deactivation Rule



- Extend both rules
- Activation: When calling a method C.foo, activate C and all of its enclosing classes
- Deactivation: Extend scope(C) such that it also includes the scope of all nested classes of C

Example 2: Scope of Classes



Modifications in Object are globally visible



Example 2: Invocation Code





Object {Object, AddressBook, AddressBook.Address, Networking, Networking, Address, String}

AddressBook			
- Address ^{Address} to_address	– Address └ ^{Address} └ to_address		
L String	– String		
to_address	└ to_address		
	Pinging ^(Pinging) ping		



S = { Object }

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Object {Object, AddressBook, AddressBook.Address, Networking, Networking.Address, String}

AddressBook				
– Address ^{Address} to_address	- Address └ ^{Address} to_address			
L String	– String			
to_address	to_address			
└─ Pinging └ ^(Pinging) └─ping				

S = { Object, Application }





Object {Object, AddressBook, AddressBook.Address, Networking, Networking, Address, String}

- AddressBook (AddressBook, Address, String) (Networking, Address, String) (String)				
– Address	– Address			
_ String	– String			
to_address	└ to_address			
	Pinging ^{Pinging} ping			

S = { Object, Networking, Pinging }







Implementation



- Prototypical implementation using metaprogramming
- Uses debug_inspector API for stack walking to implement customized method lookup in Ruby
- Give it a try (use Ruby 2.3): git@github.com:matthias-springer/ruby-class-ext.git

Related Work



- Classboxes ^[Bergel03]: Additional organizational unit (classbox), no support class nesting hierarchies
- Ruby Refinements:

Pure lexical scoping (no local rebinding)

- Context-oriented Programming (COP) ^[Hirschfeld08]: Manual activation/deactivation necessary, difficult to control when modifications should be deactivated
- *Method Shells* ^[Takeshita13]: Additional organizational unit (method shell), new syntax for including/linking
- MultiJava ^[Clifton00], Expanders ^[Warth06]: No dynamic scoping, no new methods

Summary



"Extension Classes":

- A new approach for open classes in Ruby
- Avoiding destructive modifications
- Reusable modifications (via modules)
- Scoped with respect to class nesting hierarchies
- Classes as only organizational unit

References



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Appendix

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Ruby Refinements



module NoPopup refine WebPage do def popup; end end end class Browser using NoPopup def open(url) WebPage.new.open(url) end Pure lexical scoping: NoPopup will be end deactivated after calling WebPage.open



Combining Modifications



- What happens if multiple classes with variation points for the same method are active?
- S is actually not a set but a stack
 →Class composition stack
 (cf. layer composition stack in COP ^[Hirschfeld08])
- Last activated class takes precedence
- Modified super keyword to navigate 3 hierarchies
 - Inheritance hierarchy of layer class (i.e., of class containing modifications) → takes care of mixins
 - Class composition stack (proceed in COP, AOP)
 - Inheritance hierarchy of receiver class

Method Lookup

- 1. Superclass inheritance hierarchy of layer class
- 2. Layer composition stack
- 3. Receiver class inheritance hierarchy



superclass hierarchy of layer

active layer



Example 3: Overview



Object {Object, AST, Nodes, Node, IntNode, PlusNode, Application}



Example 3: Mixins



- Conceptually, module inclusion (mixin application) creates a new superclass
- Formalism does not have a special rule for mixin application, but only for superclasses
 - → Assume that modules have been *desugared* to explicit superclasses from now on

Example 3: Inheritance





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Ex. 3: Method Lookup for IntNode





explicit superclasses

[...]: Partial class (conceptually)

Assuming single active class Application-

- Activation/deactivation rules remain unchanged
- *Effective superclass hierarchy* determines method to be executed (and also guides the lookup for proceed (super))

- Application [IntNode]
- Evaluating() [IntNode]
- 3. Printing() [IntNode]
- 4. Object [IntNode]
- 5. IntNode

1.

2.

9.

- 6. Application [Node]
- 7. Evaluating() [Node]
- 8. Printing() [Node]
 - Object [Node]
- 10. Node
- 11. Application [Object]
- 12. Evaluating() [Object]
- 13. Printing() [Object]
- 14. Object [Object]
- 15. Object

Def.: Effective Superclass Hierarchy

Definition. The effective superclass hierarchy of a class C is defined as Effective(C), where S is the class composition stack (S[1] is top of stack), #C is the number of superclasses of a class C, $super^{i}(C)$ is the i-th superclass of class C, L[C] is the partial class targeting C defined in L (if there is one), $\langle \rangle$ brackets denote a (ordered) list, and summation is used for list concatenation.

$$LayerHierarchy(L, C) = \sum_{i=0}^{\#L} \langle super^{i}(L)[C] \rangle$$
$$ClassLayers(C) = \left(\sum_{i=1}^{|S|} LayerHierarchy(S[i], C) \right) + \langle C \rangle$$
$$Effective(C) = \sum_{i=0}^{\#C} ClassLayers(super^{i}(C))$$

LayerHierarchy(L, C) is the list of partial classes for class C defined in class L and its superclasses. ClassLayers(C) is the list of partial classes of C (among all activated classes) and C itself. S = (Application)

In Example 3:

LayerHierarchy(Application, IntNode) = (Application [IntNode], Evaluating() [IntNode], Printing() [IntNode], Object [IntNode])

Same as above plus IntNode (only one layer/active class, i.e., Application in this example)

Account for superclasses of IntNode



Definition. The scope of a class L is defined as the set containing L



(reflexivity)

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Definition. The scope of a class L is defined as the set containing L, all target classes (and their reachable nested classes^{*6}) corresponding to partial classes of L, all classes in the scope of all nested classes of L^{*7} .





Definition. The scope of a class L is defined as the set containing L, all target classes (and their reachable nested classes^{*6}) corresponding to partial classes of L, all classes in the scope of all nested classes of L^{*7} , and all classes in the scope of the superclass of L

$$scope(L) = \{L\}$$
 (reflexivity)

 $\bigcup \{C \mid C \in nested^*(target(P)) \land P \in partials(L)\}$ $(dynamic \ sconing + local \ rebinding \ (+ \ hierarch \ sconing))$ $\forall isible \ in \ nested \ classes + target \ classes \ of \ superclasses \\ \cup \{C \mid C \in scope(V) \ merch \ scoping)$ $\forall scope(superclass(L))$ $(inheritance \ scoping)$



Definition. The scope of a class L is defined as the set containing L, all target classes (and their reachable nested classes^{*6}) corresponding to partial classes of L, all classes in the scope of all nested classes of L^{*7} , and all classes in the scope of the superclass of L (if super(L) \neq Object).

 $scope(L) = \begin{cases} \text{Special rule because Object is the superclass of all} \\ \text{classes. scope(Object) contains all classes.} \\ \cup \{C \mid C \in nested^{(target(P))} \land P \in partials(L) \} \end{cases}$

(dynamic scoping + local rebinding (+ hierarch. scoping))

 $\cup \{C \mid C \in scope(N) \land N \in nested(L)\}$ (hierarch. scoping)

 \cup scope(superclass(L))

(inheritance scoping)